

## CS 370 Final Definition, Theorem and Examples

**FPNS**  $(t, \beta, L, U) \pm 0.d_1 d_2 \dots d_t \times \beta^p$  where  $d_1 \neq 0$ ,  $d_i \in \{0, \dots, \beta - 1\}$ . IEEE Single Precision: 1, 8, 23, -126, 127. IEEE Double Precision: 1, 11, 52, -1022, 1023.

**AbsErr**  $= |x - \bar{x}|$ . **RelErr**  $= \frac{|x - \bar{x}|}{x}$ . **MaxRelErr**: **Machine Epsilon**. The smallest number that when added to 1 becomes the next number  $\beta/2 \cdot \beta^{-t}$ .  $\frac{\mu(x) - x}{x} = \delta$  where  $|\delta| \leq E$ .  $\mu(x) = (1 + \delta)x$ . If not in range (overflow | underflow)  $P' = P + \frac{1}{R} e d^T$ . **Markov Transition Matrix**  $M = \alpha P' + (1 - \alpha) \frac{1}{R} e e^T$ . One click:  $p^{n+1} = M p^n$ .  $p = M p$ .  $(I - M)p = 0$ .

**Eigenvector**  $x$ :  $Qx = \lambda x$ . Since  $p = M p$ , hence has eigenvector 1.

Solving upper triangular matrix: back substitution. Solving lower triangular matrix: forward substitution. Gaussian Elimination.

**LU factorization**: any square matrix  $A$  can be factored into a product of an upper triangular and lower triangular matrices such that  $LU = PA$ .

LU factorization Applications

1. Solving  $Ax = b$ .  $Ax = b \Rightarrow PAx = Pb \Rightarrow LUx = Pb$ . Solve  $Lz = Pb$ . Solve  $Ux = z$ . LU factorization  $O(N^3)$  + forward sub + back sub  $O(N^2)$ .

2. Solving  $AX = B$ . LU factorization  $O(N^3)$ . Solve  $LUx_i = Pb_i$  which is  $O(N^2)$  each.

LU factorization Example: **note** put the one that has greatest absolute value on top.

$$\begin{bmatrix} 1 & 1 & 3 & 0 \\ 2 & 1 & -1 & 1 \\ -1 & 2 & 3 & -1 \\ 3 & -1 & -1 & 2 \end{bmatrix} \xrightarrow{\text{Swap } 1,4} \begin{bmatrix} 3 & -1 & -1 & 2 \\ 2 & 1 & -1 & 1 \\ -1 & 2 & 3 & -1 \\ 1 & 1 & 3 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} 3 & -1 & -1 & 2 \\ 0 & \frac{5}{3} & -\frac{1}{3} & -\frac{1}{3} \\ 0 & \frac{5}{3} & \frac{8}{3} & -\frac{1}{3} \\ 0 & \frac{4}{3} & \frac{10}{3} & -\frac{2}{3} \end{bmatrix} \rightarrow \begin{bmatrix} 3 & -1 & -1 & 2 \\ 0 & \frac{5}{3} & -\frac{1}{3} & -\frac{1}{3} \\ 0 & 0 & 3 & 0 \\ 0 & 0 & \frac{18}{5} & -\frac{2}{5} \end{bmatrix} \xrightarrow{\text{Swap } 3,4} \begin{bmatrix} 3 & -1 & -1 & 2 \\ 0 & \frac{5}{3} & -\frac{1}{3} & -\frac{1}{3} \\ 0 & 0 & \frac{18}{5} & -\frac{2}{5} \\ 0 & 0 & 3 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} 3 & -1 & -1 & 2 \\ 0 & \frac{5}{3} & -\frac{1}{3} & -\frac{1}{3} \\ 0 & 0 & \frac{18}{5} & -\frac{2}{5} \\ 0 & 0 & 0 & \frac{1}{3} \end{bmatrix}, L = \begin{bmatrix} 1 & 0 & 0 & 0 \\ \frac{2}{3} & 1 & 0 & 0 \\ \frac{1}{3} & \frac{4}{5} & 1 & 0 \\ -\frac{1}{3} & 1 & \frac{5}{6} & 1 \end{bmatrix}, U = \begin{bmatrix} 3 & -1 & -1 & 2 \\ 0 & \frac{5}{3} & -\frac{1}{3} & -\frac{1}{3} \\ 0 & 0 & \frac{18}{5} & -\frac{2}{5} \\ 0 & 0 & 0 & \frac{1}{3} \end{bmatrix}, P = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}.$$

**Singular Value Decomposition**:  $A = U \Sigma V$ .  $U, V$  are unitary matrices. Unitary matrix:  $M^T = M^{-1}$ .

SVD Applications

1. The rank of  $A$  is the number of nonzero singular values

2. Lower-Rank Approximation using the SVD. If you want a rank- $k$  approximation to  $A$ , simply set all singular values to zero except for the largest  $k$ . (used for image compression)

3. query matching

**Numerical Solution of ODEs**.  $\frac{dy(t)}{dt} = f(t, y)$ ,  $y(t_0) = y^{(i)}$ .

**Euler's Method**. Starting with the initial state, we can approximate the solution by taking small steps in time.  $y_{k+1} = y_k + h_k f(t_k, y_k)$ . Local error  $l_{n+1} = |\hat{y}_n(x_{n+1}) - y_{n+1}|$ . **Global error**  $\Sigma_{n+1} = |\hat{y}_0(x_{n+1}) - y_{n+1}|$ . The **local error** is  $O(h^2)$ , the global error is  $O(h)$ .

Euler's Method Golf Example:  $(x(0), y(0)) = (0, 0)$ . Golfer hit the ball with initial velocity vector  $(v_x, v_y)$ . Dynamics Model:

$\frac{dx(t)}{dt} = v_x$ ,  $\frac{d^2y(t)}{dt} = -g$ . Let  $z_1 = x$ ,  $z_2 = y$ ,  $z_3 = \frac{dy}{dt} = \frac{dz_2}{dt}$ . Then

$$\frac{d}{dt} \begin{bmatrix} z_1 \\ z_2 \\ z_3 \end{bmatrix} = f(t, z) = \begin{bmatrix} v_x \\ z_3 \\ -g \end{bmatrix} \quad \& \quad z_0 = \begin{bmatrix} 0 \\ 0 \\ v_y \end{bmatrix}$$

Let  $v_x = 30$ ,  $v_y = 12$ ,  $g = 9.81$ ,  $h = 0.1$ . Then

$$z^{(1)} = z^{(0)} + h f(t_0, z^{(0)}) = \begin{bmatrix} 0 \\ 0 \\ 12 \end{bmatrix} + 0.1 \begin{bmatrix} 30 \\ 12 \\ -9.81 \end{bmatrix} = \begin{bmatrix} 3 \\ 1.2 \\ 11.019 \end{bmatrix}$$

$$z^{(2)} = z^{(1)} + h f(t_1, z^{(1)}) = \begin{bmatrix} 3 \\ 1.2 \\ 11.019 \end{bmatrix} + 0.1 \begin{bmatrix} 30 \\ 11.019 \\ -9.81 \end{bmatrix} = \begin{bmatrix} 6 \\ 2.3019 \\ 10.038 \end{bmatrix}$$

**Modified Euler Method** (improved euler OR 2nd-order Runge-kutta)

1. Start with an euler step.  $y_{n+1}^E = y_n + h_n f(t_n, y_n)$ . Let  $\bar{f}_1 = f(t_n, y_n)$ .

2.  $\bar{f}_2 = f(t_{n+1}, y_{n+1}^E)$ .

3.  $y_{n+1}^M = y_n + h_n \frac{\bar{f}_1 + \bar{f}_2}{2}$ .

The local error is  $O(h^3)$ . Global error is  $O(h^2)$ . Use Taylor expansion to **calculate local error**.

$$P' = \begin{bmatrix} 0 & 0 & 0 & \frac{1}{4} \\ 1 & 0 & \frac{1}{2} & \frac{1}{4} \\ 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 0 & \frac{1}{2} & \frac{1}{4} \end{bmatrix}$$

Given  $n$  data points  $(x_i, y_i)$ ,  $i = 1, \dots, n$  with  $x_i \neq x_j$  if  $i \neq j$ ,  $\exists!$  polynomial of degree at most  $n - 1$ ,  $p(x) = c_1 + c_2x + \dots + c_nx^{n-1}$  such that  $p(x_i) = y_i$ .

**Monomial Form:** Vandermonde System:  $n$  equations,  $n$  unknowns. **Lagrange Form:**  $L_i(x) = \frac{\dots(x - x_{i-1})(x - x_{i+1})\dots}{\dots(x_i - x_{i-1})(x_i - x_{i+1})\dots}$ .

$$p(x) = \sum_{i=1}^n L_i(x)y_i.$$

**Monomial drawback:** solve a linear system, matrix entries gets larger as  $n$  gets bigger, difficult to invert as you include more points, especially if two points have similar  $x$ -values. Difficult to solve accurately.

Piecewise (polynomial, linear) interpolation; **Cubic spline interpolation:**  $S(x)$  is called a cubic spline if 1).  $S(x)$  is an interpolant. 2).  $S(x)$  is piecewise cubic. 3).  $S(x)$  is twice differentiable.

There are  $4n - 4$  unknowns for all cubic pieces. There are  $2n - 2$  equations for interpolant,  $n - 2$  for continuous of  $S'(x)$ ,  $n - 2$  for continuous of  $S''(x)$ . The last 2 equations can be 1). Clamped spline  $S'(x_1)$  and  $S'(x_n)$  are specified. 2). Natural spline  $S''(x_1) = S''(x_n) = 0$ . 3). Periodic spline  $S'(x_1) = S'(x_n)$ ,  $S''(x_1) = S''(x_n)$  (assumes  $y_1 = y_n$ ). 4). Combinations

Use  $p_i(x) = a_{i-1} \frac{(x_{i+1} - x)^3}{6h_i} + a_i \frac{(x - x_i)^3}{6h_i} + b_i(x_{i+1} - x) + c_i(x - x_{i-1})$ . Interpolant constant gives  $b_i = \frac{y_i}{h_i} - a_{i-1} \frac{h_i}{6}$ ,  $c_i = \frac{y_{i+1}}{h_i} - a_i \frac{h_i}{6}$ .  $S''(x)$  is continuous by design.

$$p'_i(x_{i+1}) = p'_{i+1}(x_{i+1}) \Rightarrow a_{i-1} \frac{h_i}{6} + a_i \frac{h_i + h_{i+1}}{3} + a_{i+1} \frac{h_{i+1}}{6} = \frac{y_{i+2} - y_{i+1}}{h_{i+1}} - \frac{y_{i+1} - y_i}{h_i}$$

$$\begin{bmatrix} \frac{h_1}{3} & \frac{h_1}{6} & \dots & & & \\ \frac{h_1}{6} & \frac{h_1+h_2}{3} & \frac{h_2}{6} & & & \\ & \ddots & \ddots & \ddots & & \\ & & \frac{h_i}{6} & \frac{h_i+h_{i+1}}{3} & \frac{h_{i+1}}{6} & \\ & & & \ddots & \ddots & \ddots \\ & & & & \frac{h_{n-1}}{6} & \frac{h_{n-1}}{3} \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ \vdots \\ a_i \\ \vdots \\ a_{n-1} \end{bmatrix} = \begin{bmatrix} \frac{y_2-y_1}{h_1} - v_1 \\ \frac{y_3-y_2}{h_2} - \frac{y_2-y_1}{h_1} \\ \vdots \\ \frac{y_{i+2}-y_{i+1}}{h_{i+1}} - \frac{y_{i+1}-y_i}{h_i} \\ \vdots \\ v_2 - \frac{y_n-y_{n-1}}{h_{n-1}} \end{bmatrix} \quad \text{where } v_1, v_2 \text{ are slope at end points}$$

**Bezier curve** 1). passes through  $(x_0, y_0), (x_N, y_N)$ , 2). tangent to  $point_1 \rightarrow point_2$ , 3). tangent to  $point_{N-1} \rightarrow point_N$

**Fourier series**  $f(x) = a_0 + \sum_{k=1}^{\infty} [a_k \cos(\frac{2\pi kx}{N}) + b_k \sin(\frac{2\pi kx}{N})] = \sum_{k=-\infty}^{\infty} C_k (\cos(\frac{2\pi kx}{N}) + i \sin(\frac{2\pi kx}{N}))$ , with  $C_k = \overline{C_{-k}}$ .

**Complex Inner Product:**  $\langle a, b \rangle = \sum_{n=1}^N a_n \overline{b_n} = \overline{\langle b, a \rangle}$ .

$$W_N^k = (e^{2\pi i/N})^k = e^{2\pi i k/N}. f_n = \frac{1}{N} \sum_{k=0}^{N-1} F_k W_N^{nk} = \frac{1}{N} \langle \vec{F}, \overline{W_N(n)} \rangle, F_k = \sum_{n=0}^{N-1} f_n \overline{W_N^{nk}} = \langle \vec{F}, W_N(k) \rangle.$$

**Discrete Fourier transform**  $M = \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & \overline{W}^1 & \overline{W}^2 & \dots & \overline{W}^{N-1} \\ \vdots & & & & \\ 1 & \overline{W}^{N-1} & \overline{W}^{(N-1)2} & \dots & \overline{W}^{(N-1)(N-1)} \end{bmatrix}, \vec{F} = M\vec{f}, \vec{f} = \frac{1}{N} \overline{M}\vec{F}. M^{-1} = \frac{1}{N} \overline{M},$

where  $\overline{M}$  is the matrix formed by conjugating everything in  $M$ .

$$\begin{bmatrix} -1 \\ 2 \\ 4 \\ 3 \\ 1 \\ 3 \\ 4 \\ 2 \end{bmatrix} \Rightarrow \begin{bmatrix} 0 \\ 5 \\ 8 \\ 5 \\ \frac{-2(W_8^0)}{-\frac{\sqrt{2}}{2} + \frac{\sqrt{2}}{2}i(\overline{W_8^1})} \\ 0(W_8^2) \\ -\frac{\sqrt{2}}{2} - \frac{\sqrt{2}}{2}i(\overline{W_8^3}) \end{bmatrix} \Rightarrow \begin{bmatrix} 8 \\ 10 \\ -8 \\ 0 \\ -2 \\ -\sqrt{2} \\ -2 \\ \sqrt{2} \end{bmatrix} \Rightarrow \begin{bmatrix} 18 \\ -2 \\ -8 \\ -8 \\ -2 - \sqrt{2} \\ -2 + \sqrt{2} \\ -2 + \sqrt{2} \\ -2 - \sqrt{2} \end{bmatrix} \Rightarrow \begin{bmatrix} 18 \\ -2 \\ -8 \\ -8 \\ -2 - \sqrt{2} \\ -2 + \sqrt{2} \\ -2 + \sqrt{2} \\ -2 - \sqrt{2} \end{bmatrix} \Rightarrow \begin{bmatrix} 18 \\ -8 \\ -2 \\ -8 \\ -2 - \sqrt{2} \\ -2 + \sqrt{2} \\ -2 + \sqrt{2} \\ -2 - \sqrt{2} \end{bmatrix} \Rightarrow \begin{bmatrix} 18 \\ -2 - \sqrt{2} \\ -8 \\ -2 + \sqrt{2} \\ -2 \\ -2 + \sqrt{2} \\ -8 \\ -2 - \sqrt{2} \end{bmatrix}$$

**Least square**  $y = A\beta + r$ ,  $E(\beta) = r^T r = \|r\|^2$ , differentiate and set to zero gives  $A^T A\beta = A^T y$ , to solve, multiply by  $(A^T A)^{-1}$ ,  $\beta = (A^T A)^{-1} A^T y$ .